

Positionality of Dumont-Thomas numeration systems

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Definition

Numeration system over the domain $\mathbb{D} \in \{\mathbb{N}, \mathbb{Z}\}$:

- *representation map* $\text{rep}: \mathbb{D} \rightarrow A^*$
- *evaluation map* $\text{val}: L \rightarrow \mathbb{D}$, where $\text{rep}(\mathbb{D}) \subseteq L \subseteq A^*$

such that $\text{val} \circ \text{rep} = \text{id}_{\mathbb{D}}$.

Common ways of defining numeration systems:

- *Positional* numeration systems: given $(U_n)_{n \in \mathbb{N}}$, set $\text{val}_U: A^* \rightarrow \mathbb{N} : w_{\ell-1} \cdots w_0 \mapsto \sum_{i=0}^{\ell-1} w_i U_i$.
- *Abstract* numeration systems: if A is ordered, given a regular language $L \subseteq A^*$, val_L is the unique increasing bijection between $(L, \preccurlyeq_{\text{rad}})$ and $(\mathbb{N}, \leq)$.

Positionality (1)

Abstract numeration systems may or may not be also definable in a positional way.

Example

The abstract numeration system defined from $L = \{0, \dots, 9\}^* \setminus 0\{0, \dots, 9\}^*$ is equal to the usual decimal system, corresponding to $U_i = 10^i$.

Example

The abstract numeration system defined from $L = 1^*2^*$ verifies $\text{rep}(3) = 11$ and $\text{rep}(5) = 22$, so it cannot be defined as a positional numeration system.

Positionality (2)

Definition

A numeration system over $\mathbb{N}$ is *positional* if the underlying alphabet A is a set of consecutive integers $\{0, 1, \dots, c\}$ for some $c \in \mathbb{N}$ and if there exists a sequence $(U_i)_{i \geq 0} \in \mathbb{N}^{\mathbb{N}}$ such that the evaluation map is of the form

$$\text{val}: A^* \rightarrow \mathbb{N}, w_{k-1} \cdots w_0 \mapsto \sum_{i=0}^{k-1} w_i U_i.$$

Positionality (2)

Definition

A numeration system over $\mathbb{Z}$ is *positional* if the underlying alphabet A is a set of consecutive integers $\{0, 1, \dots, c\}$ for some $c \in \mathbb{N}$ and if there exist sequences $(U_i)_{i \geq 0}, (V_i)_{i \geq 0} \in \mathbb{N}^{\mathbb{N}}$ such that the evaluation map is of the form

$$\text{val}: A^* \rightarrow \mathbb{Z}, w_{k-1} \cdots w_0 \mapsto \sum_{i=0}^{k-2} w_i U_i - w_{k-1} V_{k-1}.$$

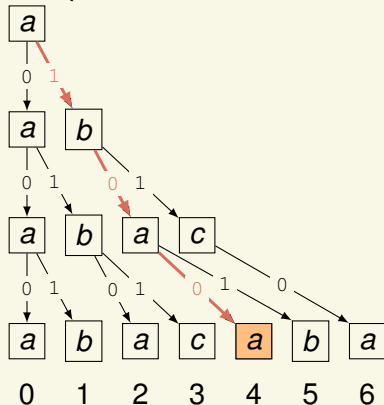
Example

The two's complement numeration system is defined by the weights $U_n = V_n = 2^n$. For instance, we have

$$\text{val}(1011) = -8 + 2 + 1 = -5.$$

Dumont-Thomas numeration systems (1989)

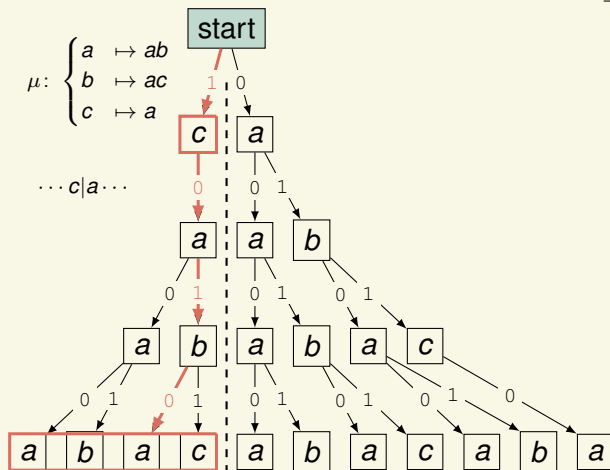
$$\mu: \begin{cases} a \mapsto ab \\ b \mapsto ac \\ c \mapsto a \end{cases}, \text{ fixed point starting with } a.$$



Use the label of a shortest path from the root to column n to represent n : $\text{rep}(4) = 100$.

Generalized Dumont-Thomas numeration systems(1)

Motivation: use Dumont-Thomas numeration systems as a way to generate Wang tilings $\rightsquigarrow$ we must be able to deal with negative numbers $\rightsquigarrow$ we must be able to deal with periodic points of substitutions.



n	$\text{rep}(n)$
7	0001000
6	0110
5	0101
4	0100
3	0011
2	0010
1	0001
0	0
-1	1
-2	1010
-3	1001
-4	1000
-5	1010110

Generalized Dumont-Thomas numerations (2)

Definition (Labbé, Lepšová, 2024; K., Labbé, Stipulanti, 2025)

Given a substitution μ , a periodic point u of period p of μ with seed $b|a$ and some $r \in \{0, \dots, p-1\}$, the *complement Dumont-Thomas numeration system associated with μ , u and r* is defined by its representation map, such that $\text{rep}(n)$ is the label of a shortest path of length congruent to $r \bmod p$, going from the root to a node in column n in the tree $\mathcal{T}_{\mu, b|a}$.

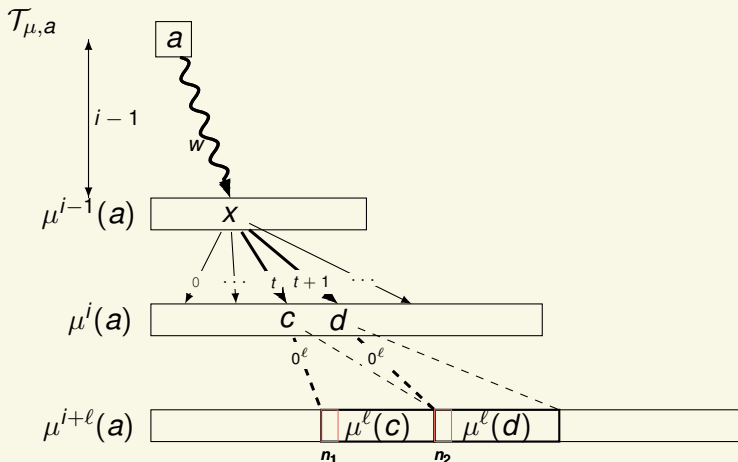
Question

Given μ, u, r , can we decide the positionality of the associated Dumont-Thomas numeration system?

If we consider $\mu: a \mapsto aab, b \mapsto a$ and $\rho: a \mapsto abb, b \mapsto ab$, both with seed $b|a$, the system associated with μ is positional but the one associated with ρ is not.

Sketch of the argument (1)

Assume the system is positional and let us compare the values of the words $wt0^\ell$ and $w(t+1)0^\ell$.



$\text{val}(wt0^\ell) = n_1$, $\text{val}(w(t+1)0^\ell) = n_2$, so $U_\ell = n_2 - n_1 = |\mu^\ell(c)|$.

Sketch of the argument (2)

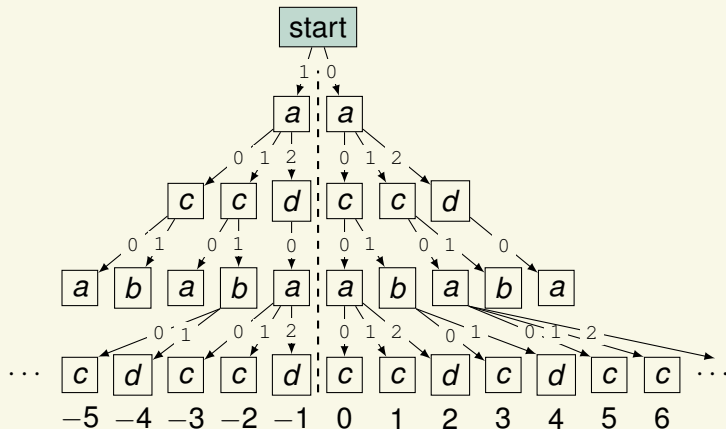
Conjecture (K., Labbé, Stipulanti, 2024)

If the Dumont–Thomas numeration system associated with μ , u , and r is positional, then $|\mu^\ell(c)| = U_\ell$ for every letter c that has a younger sibling in $\mathcal{T}_{\mu,b|a}$.

Other direction: If $|\mu^\ell(c)| = U_\ell$ for every letter c that has a younger sibling in $\mathcal{T}_{\mu,b|a}$, then incrementing the digit at position ℓ in an expansion increases the value by U_ℓ , so the system is indeed positional.

Complications (1)

$\mu: a \mapsto ccd, b \mapsto cd, c \mapsto ab, d \mapsto a$, seed $a|a$.

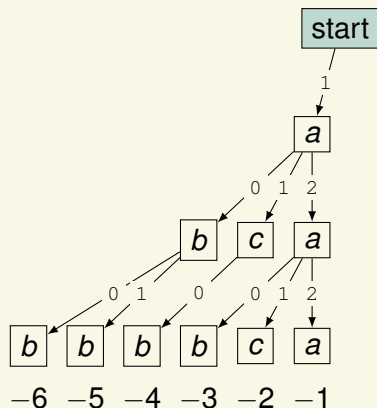


The system is positional for both values of r , despite the fact that $|\mu^\ell(a)| \neq |\mu^\ell(c)|$ for all odd ℓ :

ℓ	0	1	2	3	4
$ \mu^\ell(a) $	1	3	5	13	21
$ \mu^\ell(c) $	1	2	5	8	21

Complications (2)

$\mu: a \mapsto bca, b \mapsto bb, c \mapsto b$ with the seed $a|b$.



n	$\text{rep}_{u,0}(n)$
-1	1
-2	11
-3	10
-4	110
-5	101
-6	100
-7	1101
weight	...421

Positional ($U_\ell = 2^\ell$, $V_0 = 1$, $V_\ell = 3.2^\ell$), despite the fact that $|\mu^\ell(b)| \neq |\mu^\ell(c)|$ for all $\ell \geq 1$.

Complications (3)

In the first example, the letters a and c occur only at levels of a given parity in the tree, so the sketch can only be applied for half the values of ℓ .

In the second example, the letter to the right of a c can never be part of a shortest path to a column, so the sketch cannot be applied to c .

Definition (Technical)

For $j \in \{0, \dots, p-1\}$, the set E_j is the set of letters c verifying one of the following conditions:

- There exists a node in $\mathcal{T}_{\mu, b|a}$ labeled by c at a level congruent to $j \bmod p$, in a column other than -2 , that has a younger sibling.
- The node at level j and column -2 in $\mathcal{T}_{\mu, b|a}$ is labeled by c , it has a younger sibling labeled by d , and $|\mu^{p-j}(d)| > 1$.

We let μ be a substitution, u a periodic point with period p and seed $b|a$, and $r \in \{0, \dots, p-1\}$.

Theorem (K., Labbé, Stipulanti 2025)

The Dumont-Thomas numeration system associated with μ, u and r is positional if and only if both of the following occur:

- The map $c \mapsto |\mu^\ell(c)|$ has a constant value U_ℓ over E_j for all ℓ, j such that $\ell + j \equiv r \pmod{p}$.
- (Technical) For $j \in \{0, \dots, r-1\}$, if the node at level j and column -2 in $\mathcal{T}_{\mu, b|a}$ is labeled by c and has a younger sibling labeled by d with $|\mu^{p-j}(d)| = 1$, then $|\mu^{r-j}(c)| = U_{r-j}$.

In this case, $(U_\ell)_\ell$ and $V_\ell = |\mu^\ell(b)|$ are the sequences of weights of the system.

Corollaries (1)

In some cases, the complications outlined above do not occur. The first is the case of a primitive substitution.

We say that a letter is *non-final* in the substitution μ if it occurs in the image of any letter at any position other than the last one.

Corollary

Let μ be a primitive substitution, u be a periodic point of μ with seed $b|a$ and period p , and $r \in \{0, \dots, p-1\}$. The Dumont-Thomas complement system associated with μ , u and r is positional if and only if the map $c \mapsto |\mu^\ell(c)|$ is constant over the non-final letters in μ for every ℓ .

In this case, the value of the constant is U_ℓ , and $V_\ell = |\mu^\ell(b)|$.

Corollaries (2)

The second case where no complication occurs is that of the original Dumont-Thomas numeration systems.

Corollary

Let μ be a substitution and u be a right-infinite fixed point of μ . The Dumont-Thomas numeration system associated with μ and u is positional if and only if the map $c \mapsto |\mu^\ell(c)|$ is constant over the non-final letters in μ for every ℓ . In this case, the value of the constant is U_ℓ .

Example

Consider the substitution $\mu: a \mapsto abc, b \mapsto aac, c \mapsto a$. The non-final letters are a and b , and we can show by induction that their images by μ^ℓ have the same length for any ℓ . Thus, the system is positional.

Bertrand numeration systems

Bertrand numeration systems (Bertrand-Mathis, 1989; Charlier, Cisternino, Stipulanti, 2022) are a special case of greedy positional numeration systems, defined by either of the two properties:

- $w \in \text{rep}(\mathbb{N}) \Leftrightarrow w0 \in \text{rep}(\mathbb{N})$ for any nonempty w .
- The lexicographically largest words of each length in $\text{rep}(\mathbb{N})$ are all prefixes of one another.

There are three kinds of Bertrand numeration systems:

- $U_\ell = \ell + 1$ (trivial).
- $U_\ell = d_1 U_{\ell-1} + d_2 U_{\ell-2} + \dots + d_\ell U_0 + 1$ where $d_1 d_2 \dots$ is the quasi-greedy Rényi representation of 1 in some base β (canonical).
- $U_\ell = d_1 U_{\ell-1} + d_2 U_{\ell-2} + \dots + d_\ell U_0 + 1$ where $d_1 d_2 \dots$ is the greedy Rényi representation of 1 in some simple Parry base β (non-canonical).

The original Dumont-Thomas numeration systems have the above properties.

Simplifying the morphism

Let us study the case where μ has a right-infinite fixed point u .

Lemma

Let $\mu: A^ \rightarrow A^*$ be a substitution such that the map $c \mapsto |\mu^\ell(c)|$ is constant over the non-final letters in μ for every integer $\ell \geq 0$. Then there exist an alphabet $B \subseteq A$, a substitution $\nu: B^* \rightarrow B^*$ such that ν has only one non-final letter and a coding $\phi: A \rightarrow B$ such that ν and $\phi(u)$ define the same Dumont-Thomas numeration system as μ and u .*

A substitution that has a fixed point and only one non-final letter is of the form

$$\mu: a_1 \mapsto a_1^{d_1} a_2, a_2 \mapsto a_1^{d_2} a_3, \dots, a_n \mapsto a_1^{d_n} a_k$$

for some $n \geq 1$, $1 \leq k \leq n$, $d_1 > 0$ and $d_2, \dots, d_n \geq 0$.

Fabre substitutions

Another approach to Bertrand numeration systems (Fabre, 1995): If β is a Parry number and the quasi-greedy Rényi representation of 1 is $(d_1 \cdots d_{k-1})(d_k \cdots d_n)^\omega$, define the substitution

$$\mu_\beta: 1 \mapsto 1^{d_1}2, 2 \mapsto 1^{d_2}3 \dots, n \mapsto 1^{d_n}k.$$

For instance, for β equal to the positive root of $x^3 - x^2 - x - 1 = 0$, the quasi-greedy representation of 1 is $(110)^\omega$ and we find

$$\mu_\beta: 1 \mapsto 12, 2 \mapsto 13, 3 \mapsto 1.$$

Theorem (Fabre, 1995)

$|\mu_\beta^\ell(1)|$ is the U_ℓ defined for the canonical Bertrand numeration system for β . This numeration system is the Dumont-Thomas numeration system associated with μ_β and $\mu_\beta^\infty(1)$.

Equivalence between the two systems (1)

The other two kinds of Bertrand numeration systems are also associated with Dumont-Thomas numeration systems when β is Parry.

The converse is not always true.

Example

Consider $\mu: a \mapsto aab, b \mapsto aaaa$. The associated Dumont-Thomas system is positional, with weight sequence $1, 3, 10, 32, \dots$. However, it is not equal to a Bertrand numeration system: $\text{rep}(9) = 23$, but it should be 30 .

If such a system were Bertrand associated with some β , the quasi-greedy β -representation of 1 would be $(23)^\omega$, which breaks the Parry conditions.

Equivalence between the two systems (2)

Proposition (K., Labbé, Stipulanti 2025)

Let μ be a substitution of the form

$$\mu : a_1 \mapsto a_1^{d_1} a_2, a_2 \mapsto a_1^{d_2} a_3, \dots, a_n \mapsto a_1^{d_n} a_k$$

for some $n \geq 1$, $1 \leq k \leq n$, $d_1 > 0$ and $d_2, \dots, d_n \geq 0$.

Construct the word $d_1 d_2 \cdots = d_1 \cdots d_{k-1} (d_k \cdots d_n)^\omega$. The Dumont–Thomas numeration system associated with μ and the seed a_1 is equal to a Bertrand numeration system if and only if we have $d_i d_{i+1} \cdots \preccurlyeq_{\text{lex}} d_1 d_2 \cdots$ for each $i \geq 1$.

- Dumont-Thomas numeration systems, interpreting substitutions as trees, are easily generalized to biinfinite periodic points of substitutions.
- Their positionality is completely understood.
- In the case of the original Dumont-Thomas numeration systems, the positional systems are those that correspond to substitutions "in the style of Fabre".
- Those systems are Bertrand numerations associated with Parry numbers if and only if a lexicographic condition is met.

Thank you!